0. Introduction

ELEG5481

SIGNAL PROCESSING OPTIMIZATION TECHNIQUES

0. INTRODUCTION

A General Formulation of Optimization

$$\min f_0(x)$$

s.t. $f_i(x) \le b_i, \quad i = 1, \dots, m$

$$x = [x_1, \ldots, x_n]^T \in \mathbf{R}^n$$
 optimization variables

- $f_0: \mathbf{R}^n o \mathbf{R}$ objective function
- $f_i: \mathbf{R}^n \to \mathbf{R}, i = 1, \dots, m$ constraint functions

Our aim is to find an **optimal solution** x^* that minimizes f_0 whilst satisfying all the constraints.

Why does optimization concern me?

- Optimization has found applications in a wide variety of areas such as finance, statistics, and engineering, to name a few.
- In engineering, it plays a key role in solving or handling numerous (and sometimes very hard) problems in control, circuit design, networks, signal processing, and communications.
- It is an important enough topic that we should know at least a bit about it.

Example: Diet Problem

- x_i is the quantity of food i.
- Each unit of food i has a cost of c_i .
- One unit of food j contains an amount a_{ij} of nutrient i.
- We want nutrient i to be at least equal to b_i .
- **Problem:** find the cheapest diet such that the minimum nutrient requirements are fulfilled.

This problem can be formulated as:

$$\min \sum_{i=1}^{n} c_i x_i$$

s.t.
$$\sum_{j=1}^{n} a_{ij} x_j \ge b_i, \quad i = 1, 2, \dots, m$$
$$x_i \ge 0, \quad i = 1, 2, \dots, n$$

This is a linear program.

Example: Chebychev Center

- Let a norm ball $B(x_c, r) = \{ x \mid ||x_c x||_2 \le r \}$, & a polyhedron $\mathcal{P} = \{ x \mid a_i^T x \le b_i, i = 1, ..., m \}.$
- **Problem:** Find the largest ball inside a polyhedron \mathcal{P} ; i.e., $\max_{x_c,r} r$, subject to $B(x_c,r) \subseteq \mathcal{P}$.



Example: Optimal Power Assignment in Wireless Communications

• Consider a wireless comm. system with K transmitters & K receivers.



• Receiver *i* is intended to receive information only from Transmitter *i*, & it sees the other transmitters as interferers.

• The signal-to-interference-and-noise ratio (SINR) at receiver *i*

$$\gamma_i = \frac{G_{ii}p_i}{\sum_{j \neq i} G_{ij}p_j + \sigma_i^2}$$

where

 p_i is the transmitter i power,

 G_{ij} is the path gain from transmitter j to receiver i,

 σ_i^2 is the noise power at receiver *i*.

• **Problem:** Maximize the weakest SINR subject to power constraints $0 \le p_i \le p_{\max,i}$, where $p_{\max,i}$ is the max. allowable power of transmitter *i*.

$$\max_{\substack{p_i \in [0, p_{\max,i}] \\ i=1,\dots,K}} \min_{i=1,\dots,K} \frac{G_{ii}p_i}{\sum_{j \neq i} G_{ij}p_j + \sigma_i^2}$$

Example: Filter Design in Digital Signal Processing

Finite impulse response (FIR) filter

$$y[k] = \sum_{i=0}^{n-1} h_i x[k-i]$$

where

 $\boldsymbol{x}[\boldsymbol{k}]$ input sequence

 $\boldsymbol{y}[\boldsymbol{k}]$ output sequence

 h_i filter coefficients

 \boldsymbol{n} filter order or length

Frequency response:

$$H(\omega) = \sum_{i=0}^{n-1} h_i e^{-j\omega i}$$

Problem: find $h = [h_0, \ldots, h_{n-1}]^T$ so that h and/or H satisfy/optimize certain given specifications.

For example, the Chebychev design solves for

$$\min_{h \in \mathbf{R}^n} \max_{\omega \in [0,\pi]} \left| H(\omega) - H_{des}(\omega) \right|$$

given a desired freq. response $H_{des}(\omega)$. This design minimizes the worst-case absolute error between the desired and actual freq. responses.



With modern software and with proper training in convex optimization, you can programme this filter design in 30 minutes.

Example: Beamformer Design in Array Signal Processing

Uniform linear array:



Signal model:

- A1) far-field situations so that source waves are planar; &
- A2) narrowband source signals so that the received signal of one sensor is a phase shifted version of that of another.

If a source signal $s(t) \in \mathbb{C}$ comes from a direction of θ , the array output $y(t) = [y_1(t), \dots, y_P(t)]^T$ is $y(t) = a(\theta)s(t)$

Here,

$$a(\theta) = [1, e^{-j2\pi d \sin(\theta)/\lambda}, \dots, e^{-j2\pi d(P-1)\sin(\theta)/\lambda}]^T \in \mathbf{C}^P$$

is the steering vector, where λ is the signal wavelength.

Beamforming:

$$\hat{s}(t) = w^H y(t)$$

where $w \in \mathbf{C}^P$ is a beamformer weight vector.

- Let $\theta_{des} \in \left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$ be the desired direction.
- A simple beamformer is $w = a(\theta_{des})$, but it does not provide good sidelobe suppression.
- **Problem:** find a w which minimizes sidelobe energy subject to a pass response to $\theta_{\rm des}$.



Direction pattern of the conventional beamformer. $\theta_{des} = 10^{\circ}$; P = 20.

- Let $\Omega = [-\pi/2, \theta_{\ell}] \cup [\theta_u, \pi/2]$ be the sidelobe band, where $[\theta_{\ell}, \theta_u]$ is the band of interest.
- Average sidelobe energy minimization:

$$\min \int_{\Omega} |w^{H}a(\theta)|^{2} d\theta$$

s.t. $w^{H}a(\theta_{des}) = 1$

• Worst-case sidelobe energy minimization:

min
$$\max_{\theta \in \Omega} |w^H a(\theta)|^2$$

s.t. $w^H a(\theta_{des}) = 1$



Direction patterns of the two beamformer designs. $\theta_{des} = 10^{\circ}$. Sidelobe suppression is applied to directions outside $[0^{\circ}, 20^{\circ}]$. Wing-Kin Ma, Dept. Electronic Eng., The Chinese University of Hong Kong

Example: L1 Signal Reconstruction

• Consider a noiseless model

$$y = Ax$$

but with m < n (underdetermined system).

• By linear algebra, there are many (infinite) possible solutions to y = Ax

$$Ax = y \iff x = A^{\dagger}y + u, \ u \in \mathcal{N}(A)$$

where $A^{\dagger} = A^T (AA^T)^{-1}$, and $\mathcal{N}(A) = \{u \mid Au = 0\}$ is the nullspace of A.

- How can we choose x from these possible solutions?
- This is what matrix textbooks would tell us to do: Least 2-norm:

 $\begin{array}{ll} \min & \|x\|_2 \\ \text{s.t.} & Ax = y \end{array}$

The solution has a closed form, $x^* = A^{\dagger}y$ (make sense intuitively).

• Least 0-norm reconstruction:

 $\begin{array}{ll} \min & \|x\|_0 \\ \text{s.t.} & Ax = y \end{array}$

where $||x||_0$ counts the number of nonzero elements in x.

- Make sense for sparse signals; i.e., signals with many zeros.
- Can prove that if the no. of zeros in the actual x is sufficiently large compared to the no. of measurements m, then 0-norm minimization leads to the ground truth.
- $\|\cdot\|_0$ is not convex. In fact, 0-norm minimization poses a very hard problem.

• Least 1-norm reconstruction:

 $\begin{array}{ll} \min & \|x\|_1 \\ \text{s.t.} & Ax = y \end{array}$

- the 'best' convex approximation to 0-norm.
- can prove that under some assumptions, 1-norm minimization is able to approach 0-norm minimization (in some probabilistic sense).
- Currently a very hot topic (in the literature it is called compressive sensing)
- 1-norm minimization is an LP

$$\min \sum_{i=1}^{n} t_i$$
s.t. $Ax = y$

$$-t_i \le x_i \le t_i, \quad i = 1, \dots, n$$



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Solving optimization problems

- Given a general opt. problem, obtaining its optimal solution can be very hard.
- In the study of *nonlinear programming* (or nonlinear opt.), various numerical algorithms have been developed to *try to* find the optimal solution. Some well known examples are the gradient descent method, and the Newton method.
- These opt. algs. require an initial guess of the solution. Also, they may only guarantee convergence to a locally optimal solution.

- There are also a variety of approaches for nonlinear opt., such as the heuristics based approach (e.g., genetic algs., ant colony opt.), and the Monte-Carlo based approach (e.g., simulated annealing). Again, convergence to a globally optimal solution is not guaranteed.
- When an opt. problem is combinatorial (or discrete), the problem is generally much harder to solve; study complexity theory for the details.
- There are certain problem classes that can be solved effectively, though.

There are problems for which the solutions can be analytically found:

Example: Least Squares (LS)

 $\min \|Ax - b\|_2^2$

• LS has a closed form solution $x^* = (A^T A)^{-1} A^T b$.

Example: Entropy Maximization

- Let y be a r.v. drawn from $\{y_1, \ldots, y_n\}$.
- Let $p_i = \operatorname{prob}(y = y_i)$.
- **Problem:** find a distribution $\{p_i\}$ such that the entropy of y is maximized.

$$\max \sum_{i=1}^{n} p_i \log(1/p_i)$$

s.t. $p_i \ge 0, \quad i = 1, \dots, n$
$$\sum_{i=1}^{n} p_i = 1$$

• The solution is well known to be $p_i = 1/n$ for all i.

- Opt. problems that can be analytically solved are considered very special cases.
- Another problem class that can be effectively handled is that of the **convex optimization** problems.

Example: Linear Programming (LP)

min
$$c^T x$$

s.t. $a_i^T x \le b_i, \quad i = 1, \dots, m$

- no analytical formula for the solution
- efficient & reliable algorithms for finding the optimal solution exist.

Convex Optimization Problems

$$\min f_0(x)$$

s.t. $f_i(x) \le b_i, \quad i = 1, \dots, m$

in which the objective & constraint functions are convex:

$$f_i(\alpha x + (1 - \alpha)y) \le \alpha f_i(x) + (1 - \alpha)f_i(y)$$

for any x, y, and for any $\alpha \in [0, 1]$.



Conic Problems

Conic optimization is a representative class of convex optimization problems.

$$\min c^T x$$

s.t. $Ax = b$,
 $x \in \mathcal{K}$

where \mathcal{K} is a **convex cone**.

The well-known linear program is conic, where $\mathcal{K} = \{ x \in \mathbb{R}^n \mid x_i \ge 0, i = 1, ..., n \}$.

Second order cone program: $\mathcal{K} = \{ x \in \mathbb{R}^{n+1} \mid \sqrt{\sum_{i=1}^{n} x_i^2} \leq x_{n+1} \}$ is the second order cone.

Semidefinite program: $\mathcal{K} = \{ X \in \mathbb{R}^{n \times n} \mid X \text{ is positive semidefinite (PSD)} \}$ is the set of PSD matrices.



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Merits of convex optimization

- Reliable and efficient algorithms exist for *many* convex opt. problems, especially those under the conic opt. problem class.
- (Surprisingly) many problems can be converted to convex opt.

Role of convex optimization in nonconvex problems

- We can use convex opt. to approximate a nonconvex problem. (Having a hard problem does not mean you should give up).
- We can use convex opt. to build a nonconvex opt. algorithm (subopt. *per se*).
 - SQP, a general-purpose nonlinear opt. solver in MATLAB, may be seen as an algorithm that uses convex opt. to sequentially process an opt. problem.

Example: Unicast Transmit Downlink Beamforming

Problem scenario: a multi-antenna basestation sends data streams, one for each user.



• Transmit signal:

$$x(t) = \sum_{i=1}^{K} w_i s_i(t)$$

where $s_i(t) \in \mathbb{C}$ is the data stream for user i, & $w_i \in \mathbb{C}^N$ its tx beamformer vector.

• Received signal of user *i*:

$$y_i(t) = h_i^T x(t) + v_i(t)$$

= $h_i^T w_i s_i(t) + \sum_{l \neq i} h_i^T w_l s_l(t) + v_i(t)$
interference

• Let $R_i = E\{h_i^*h_i^T\}$ be the channel correlation. Assuming that only the channel corr. info. is available, the SINR of user *i* is

$$\mathsf{SINR}_i = \frac{w_i^H R_i w_i}{\sum_{l \neq i} w_l^H R_i w_l + \sigma_i^2}$$

• The problem is to design $\{w_i\}_{i=1}^k$ such that the total power is minimized whilst all users have SINRs no worse than a requirement γ :

$$\min_{\substack{w_1,\ldots,w_K \in \mathbb{C}^N \\ \text{s.t.}}} \frac{\sum_{i=1}^K \|w_i\|^2}{\sum_{l \neq i} w_l^H R_i w_l + \sigma_i^2} \ge \gamma, \quad i = 1, \ldots, K$$

• This problem is nonconvex, but can be solved optimally by a semidefinite program.

Example: Multicast Transmit Beamforming

Problem scenario: a multi-antenna basestation broadcasts common information to all users.

• Received signal model for user \boldsymbol{i}

$$y_i(t) = h_i^T x(t) + v_i(t)$$

where h_i is the MISO channel from the basestation to user i, $v_i(t)$ is noise with variance σ_i^2 .

• The transmit signal:

$$x(t) = ws(t)$$

where $s(t) \in \mathbb{C}$ is the transmitted common information, & $w \in \mathbb{C}^N$ is the transmit beamformer vector.



• The design problem is to minimize the transmit power subject to SNR constraints:

$$\min_{w \in \mathbb{C}^N} \quad \frac{\|w\|^2}{\text{s.t.}}$$
s.t.
$$\frac{\mathrm{E}\{|h_i^T w s(t)|^2\}}{\sigma_i^2} \ge \gamma, \quad i = 1, \dots, K$$

where γ is pre-specified.

• The design problem can be rewritten as a complex-valued nonconvex QCQP

min
$$||w||^2$$

s.t. $w^H A_i w \ge 1, \quad i = 1, \dots, K$

where $A_i = \mathrm{E}\{h_i^* h_i^T\} / \gamma \sigma_i^2$.

• This multicast transmit beamforming problem is NP-hard in general, but can be approximated by semidefinite programming.

Example: MAXCUT

- Input: A graph G = (V, E) with weights w_{ij} for (i, j) ∈ E. Assume w_{ij} ≥ 0 and w_{ij} = 0 if (i, j) ∉ E.
- **Goal:** Divide nodes into two parts so as to maximize the weight of the edges whose nodes are in different parts.



Let $V = \{1, 2, ..., n\}$ without loss of generality. The MAXCUT problem takes the form

$$\max \sum_{i=1}^{n} \sum_{j=i+1}^{n} w_{ij} \frac{1 - x_i x_j}{2}$$

s.t. $x_i \in \{-1, +1\}, \quad i = 1, \dots, n$

Some remarks:

- MAXCUT is a combinatorial opt. problem.
- To find the optimal solution of MAXCUT is very hard.

Example: MIMO detection

• Consider a multi-antenna communications system with n transmit & m receive antennas.



• The received signal model may be expressed as

$$y = Hs + v$$

where

- $y \in \mathbf{R}^m$ multi-receiver output vector;
- $s \in \{-1, +1\}^n$ transmitted symbols;
- $H \in \mathbf{R}^{m \times n}$ MIMO (or multi-antenna) channel;
- $v \in \mathbf{R}^n$ Gaussian noise with zero mean & covariance $\sigma^2 I$.
- Maximum-likelihood (ML) detection of s:

$$\min_{s \in \{\pm 1\}^n} \|y - Hs\|_2^2$$

The ML detection problem is no easier to solve than MAXCUT.

• Nevertheless, ML MIMO detection and MAXCUT can be very effectively approximated by **semidefinite programming**.

Example: Sensor Network Localization

- Consider a sensor network where each node can communicate with the other nodes.
- Depending on applications, a sensor node may have no information on its position.



- Let $z_i = [x_i, y_i]^T$ be the position of *i* sensor node.
- A node can communicate with a neighboring node to obtain a distance information

$$d_{ij} = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$$

• A sensor network localization problem:

$$\min_{z_i \in \mathbf{R}^2, \ i=1,\dots,m} \sum_{i,j} \left(d_{ij} - \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2} \right)^2$$

Such a problem is very difficult to solve. But it can be effectively approximated by semidefinite programming.

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SDR (ML-SNL formulation), plus a 2nd-stage solution refinement by gradient descent. The distance measurements are noisy. o: true sensor locations; \diamond : anchor locations; *: SDR solution; — : gradient descent trajectory (50 iterations). Wing-Kin Ma, Dept. Electronic Eng., The Chinese University of Hong Kong

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Gradient descent ML-SNL with a random starting point. o: true sensor locations; \diamond : anchor locations; — : gradient descent trajectory (50 iterations). Wing-Kin Ma, Dept. Electronic Eng., The Chinese University of Hong Kong

Course Outline

- Theory:
 - linear algebra and matrix analysis
 - convex sets and convex functions
 - convex optimization, duality
 - conic opt.: linear program, second-order cone program, semidefinite program
- Methods:
 - interior-point methods (a brief overview)
 - subgradient methods, first-order optimization
 - nonconvex problems: optimization strategies based on convex or tractable opt.

- Applications:
 - signal and image processing: optimization of digital filters, optimization of beamforming in sensor array processing, sparse optimization and compressive sensing
 - signal estimation: signal recovery or regression via ℓ_p norm optimization
 - pattern classification: support vector machine, large-margin nearest-neighbor classification
 - wireless communications: power allocation in wireless networks, MIMO transmit optimization, mobile positioning and sensor network localization
 - information theory: capacity optimization in Gaussian frequency-selective and MIMO channels
 - circuit: geometric program and circuit optimization
 - distributed optimization, with applications to image processing and wireless networks



- Assignments, 20%.
- A written exam., 30%, in midterm. To make sure that you understand the basics.
- Project, 50%. Let you have hands-on experience with using opt. to solve signal processing or engineering problems. Problems will be provided. You can also propose problems related to your own research (subject to my approval).



- Venue and Time: Every Wed, 19:00-22:00, ERB405
- Course website: http://dsp.ee.cuhk.edu.hk/eleg5481/
- My email: wkma@ee.cuhk.edu.hk



Textbook

S. Boyd & L. Vandenberghe, *Convex Optimization*, Cambridge Univ. Press, 2004. Available online: http://www.stanford.edu/~boyd/cvxbook/

Other references

- D. P. Bertsekas, et. al, Convex Analysis and Optimization, Athenta Scientific, 2002.
- R. Fletcher, *Practical Methods of Opt.*, John Wiley & Sons, 1987.
- http://www.stanford.edu/class/ee364/